$$6 - Chaining
Recap: - Massort's Finite Lamma (UFE):
sup $\frac{1}{n} \sum_{i=1}^{n} f(a_i)^{i} \in k^2 \Rightarrow \widehat{R}(\mathcal{F}) \in k \sqrt{\frac{2 \log 1}{n}}$
- Massort's Infinite Lemma:
sup $\frac{1}{n} \sum_{i=1}^{n} f(a_i)^{i} \in k^2 \Rightarrow \widehat{R}(\mathcal{F}) \in k \sqrt{\frac{2 \log 1}{n}}$
This lecture: (- Ruisist couring (discretesation))
2 - Weat to bound RC instead of excess met
ar supried press:
S - Couring is over \mathcal{F} directly, which requires
a distance.
Metrice $f_i, g \in \mathcal{F}$, $d(f, g) = \left(\frac{1}{n} \sum_{i=1}^{n} (f(a_i) - g(a_i))^2\right)^{i/2}$
where \mathcal{F}_i 's are date points.
Motation: Define $f = \frac{1}{n} \sum_{i=1}^{n} f(a_i)^2$, $d(f, g) = here f(a_i) - f(a_i)^T \in \mathbb{R}^n$
 $\Rightarrow \|f_i\|_{\mathcal{F}}^2 = \frac{1}{n} \sum_{i=1}^{n} f(a_i)^2$, $d(f, g) = \|f(g)\|$ $(f_i)\|_{L_{\mathcal{F}}} = \|f\|)$$$

 $perme \in Massort's Finite Lemma: <math>\hat{R}(F) \leq \sup \|f\| \sqrt{2 \log |F|}$ Recell: An E-cour of F wrt a metric d is a set $\mathcal{N}_{\varepsilon} = \{g_1, \dots, g_m\}$ s.t. $\forall f \in \mathcal{F}, \exists g \in \mathcal{N}_{\varepsilon} \text{ s.t. } d(f,g) \leq \epsilon$. Covering number of F: N(E, T, d) = min { |NE| : NE is a E-cover of F} Metric entropy of $\mathcal{F} \triangleq \log N(\epsilon, \mathcal{F}, d)$ $E_{X}: (all fnes) \qquad \mathcal{F} = \left\{ f: \mathcal{F} \rightarrow [0, 1] \right\}$ - One part at every vertical line. ×,f $- \left| \mathcal{M}_{\epsilon} \right| = \left(1 + \frac{1}{2\epsilon} \right)^{\eta}$ 4e 4e 2e $\leq \left(\frac{1}{\epsilon}\right)^{\gamma}$ $\Rightarrow \mathcal{N}(\epsilon, \mathcal{F}, d) \leq \left(\frac{1}{\epsilon}\right)^n$ $Z_{(1)}$ $Z_{(2)}$ $Z_{(2)}$ ZINI (exponential in n). Ex: (non-decreasing fres): $\mathcal{F} = \{f: \mathcal{F} \rightarrow [0, 1] non-decreasing \}$ - Only include get its that are non-decrasing. ₩ 5 2.c 2.c - Hg t N_€, we reed n-1 → $\mathcal{Z}_{(1)}$ $\mathcal{Z}_{(2)}$ $\mathcal{Z}_{(2)}$

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$$k_{1} = \# \uparrow$$
 at late z_{0} : $k_{1} + k_{2} + \dots + k_{n} = m$.
- $\#$ non-decreasing $g = \int \{ (k_{1} \dots k_{n}) : \sum_{i=1}^{n} k_{i} = m, k_{i} \in \mathbb{N} \} \}$

$$= \sum_{k_{i} + k_{k} - k_{n} = m^{m} \leq n^{1/2}$$

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$$\Rightarrow \mathcal{N}(e_{i}, \forall_{i}, d) \leq n^{1/2} \quad (poly in n).$$
Theorem (Discretizedum): Let $\exists = \{f^{2}\}$ set. $f : \not{Z} \rightarrow \mathbb{R}$.
Then for $k = \sup_{i=1}^{n} \|f\|$, we have

$$\widehat{\mathcal{R}}(\forall) \leq k_{1} \frac{2 \log \mathcal{N}(e_{i} \forall_{i} d)}{n} + e_{i}, \forall e > 0.$$

$$T = T$$

$$Remark: I is due to MiFL. If $\psi = \psi / e_{i}$

$$I is due to discription in I is ψ / e_{i}

$$= hold_{i} - e_{i}f_{i} \Rightarrow we need optimize over e.$$

$$f(\forall) = \lim_{i=1}^{n} [f(x_{i}) - f(x_{i})]^{T} \text{ and } \forall = \lim_{i=1}^{n} [\pi_{i}]^{T}.$$

$$f(\forall) = \lim_{i=1}^{n} [f(x_{i}) / z_{i}] = \mathbb{E}\left[\sup_{i=1}^{n} \sqrt{\forall_{i}} / \frac{1}{2 i} \int_{i=1}^{n} = \mathbb{E}\left[\sup_{i=1}^{n} / \frac{1}{2 i} \int_{i=1}^{n} = \mathbb{E}\left[\sup_{i=1}^{n} / \frac{1}{2 i} / \frac{1}{2 i} \int_{i=1}^{n} = \mathbb{E}\left[\sup_{i=1}^{n} / \frac{1}{2 i} /$$$$$$

$$= \sup_{i \in \mathcal{J}} \langle \overline{\tau}, f \rangle = \langle \overline{\tau}, g \rangle + \langle \overline{\tau}, f - g \rangle$$

$$= i \quad \langle \overline{\tau}, g \rangle + \| \overline{\tau} \| \| \| f - g \| \| (l_{\mathcal{J}} \subset G)$$

$$= i \quad \langle \overline{\tau}, g \rangle + i \in \mathbb{R}$$

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$$= i$$

Theorem (Dudley's thm):
$$\mathcal{F}$$
 is a family of face $f: I \rightarrow R$
 $\mathcal{R}(\mathcal{F}) \leq 12 \int \frac{l_{Q} \mathcal{M}(\epsilon_{1} \mathcal{F}_{1} \mathcal{G})}{n} d\epsilon$.
Remark 6: $1 - When \epsilon$ is large, $\mathcal{M}(\epsilon_{1} \mathcal{F}_{1} \mathcal{G}) = 1$. Not really
integrating to so for a compact \mathcal{F} .
 $2 - D$ is creation error ϵ is gave b.
 $3 - For$ the above non-decreasing frees
 $\mathcal{R}(\mathcal{F}) \leq 12 \int \sqrt{\frac{l_{Q} \mathcal{M}}{l_{Q} \mathcal{M}}} dc = 12 \sqrt{\frac{l_{Q} \mathcal{M}}{l_{Q} \mathcal{M}}} \int \frac{1}{l_{q}} d\epsilon$
 $= 2l_{1} \sqrt{\frac{l_{Q} \mathcal{M}}{l_{Q} \mathcal{M}}} = \frac{1}{2} \sqrt{\frac{l_{Q} \mathcal{M}}{l_{Q} \mathcal{M}}} + \epsilon$
 \mathcal{F} rewords, \mathcal{F} is \mathcal{M}_{ℓ} o \mathcal{F} is applied to $3_{\mathcal{M}} - g_{\ell}$
 \mathcal{F} is applied to $3_{\mathcal{M}} - g_{\ell}$
 $\mathcal{F}_{\ell} = 0$
 $\mathcal{F}_{\ell} = 0$